

Updatable Public-Key Encryption, Revisited

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Updatable KEM

KEM

$(pk, sk) \leftarrow KGen()$

$(K, c) \leftarrow Enc(pk)$

$K \leftarrow Dec(sk, c)$

✗ no forward secrecy: leaked sk leaks all previously encapsulated keys

Updatable KEM

$(pk, sk) \leftarrow KGen()$

$(K, c, pk') \leftarrow Enc(pk)$

$(K, sk') \leftarrow Dec(sk, c)$

✓ forward secrecy

✓ efficient

✗ requires coordination between senders



Forward-Secure KEM

$(pk, sk) \leftarrow KGen()$

$(K, c) \leftarrow Enc(pk)$

$(K, sk') \leftarrow Dec(sk, c)$

✓ forward secrecy

✗ not efficient

Our Contributions

| | New UKEM Application | Application Security | UKEM Security | Assumption | $ pk + ctxt $ |
|----------|----------------------|----------------------|---------------|------------|-----------------|
| [JMM'19] | 2-party messaging | "almost-optimal" | IND-CPA | DH (ROM) | |

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| [ACDT'20] | group messaging | very restricted | [JMM'19](-ish) | DH (ROM) | |

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| [DKW'21] | | | IND-CCA | DH, LWE | |



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| [HLP'22] | | | [DKW'21] | DCR (ROM) | 25.5 kB |

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| [AW'23] | | | [DKW'21]— | DH (ROM) | 0.38 kB |

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| [AW'23] | | | [DKW'21]— | DH (ROM) | 0.38 kB |
| [this work] | | | IND-CCA-Members | DH (ROM) | 0.24 kB |
| [this work] | | | IND-CCA-Joiners | DH+pairing (ROM, AGM) | 0.24 kB |

UKEM Syntax and Security

UKEM Syntax and Security

Basic Syntax [ACDT'20]

$(pk_1, sk_1) \leftarrow \text{KGen}()$

$(K, c, pk') \leftarrow \text{Enc}(pk)$

$(K, sk') \leftarrow \text{Dec}(sk, c)$



Basic IND-CPA Security [ACDT'20]

- oracles Enc, Chal
- Chal returns c^* and real or random key K^*

UKEM Syntax and Security

Basic Syntax [ACDT'20]

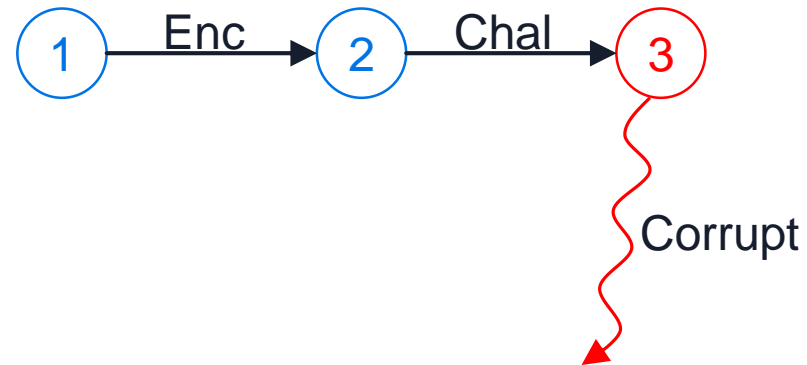
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Basic IND-CPA Security [ACDT'20]

- oracles Enc, Chal
- Chal returns c^* and real or random key K^*
- Chal also returns sk *after* decrypting c^*



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Basic IND-CCA Security [DKW'21]

add oracle $\text{Dec}(c', pk')$

UKEM Syntax and Security



Extended Syntax [DKW'21]

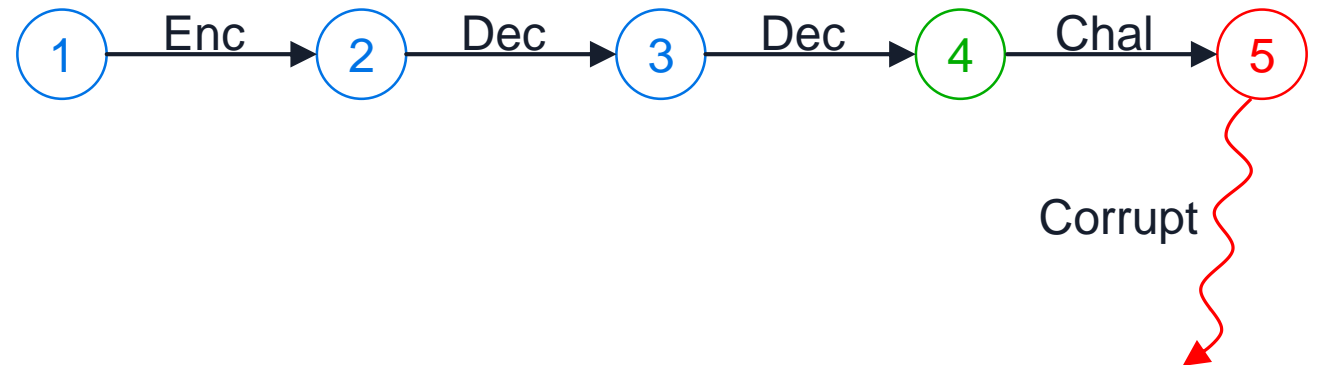
$(pk1, sk1) \leftarrow \text{KGen}()$

$(K, c, pk', mt') \leftarrow \text{Enc}(pk)$

$(K, sk') \leftarrow \text{Dec}(sk, c)$

$0/1 \leftarrow \text{Verify}(pk, pk', mt')$

allows a sender to verify a public key without the secret key



Basic IND-CCA Security [DKW'21]

add oracle $\text{Dec}(c', pk', mt')$

- if $\text{Dec}(sk, c')$ succeeds, create **full node**
- else if $\text{Verify}(pk, pk', mt')$ succeeds, create **half node**

UKEM Syntax and Security



Extended Syntax [DKW'21]

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$(K, c, pk', mt') \leftarrow Enc(pk)$

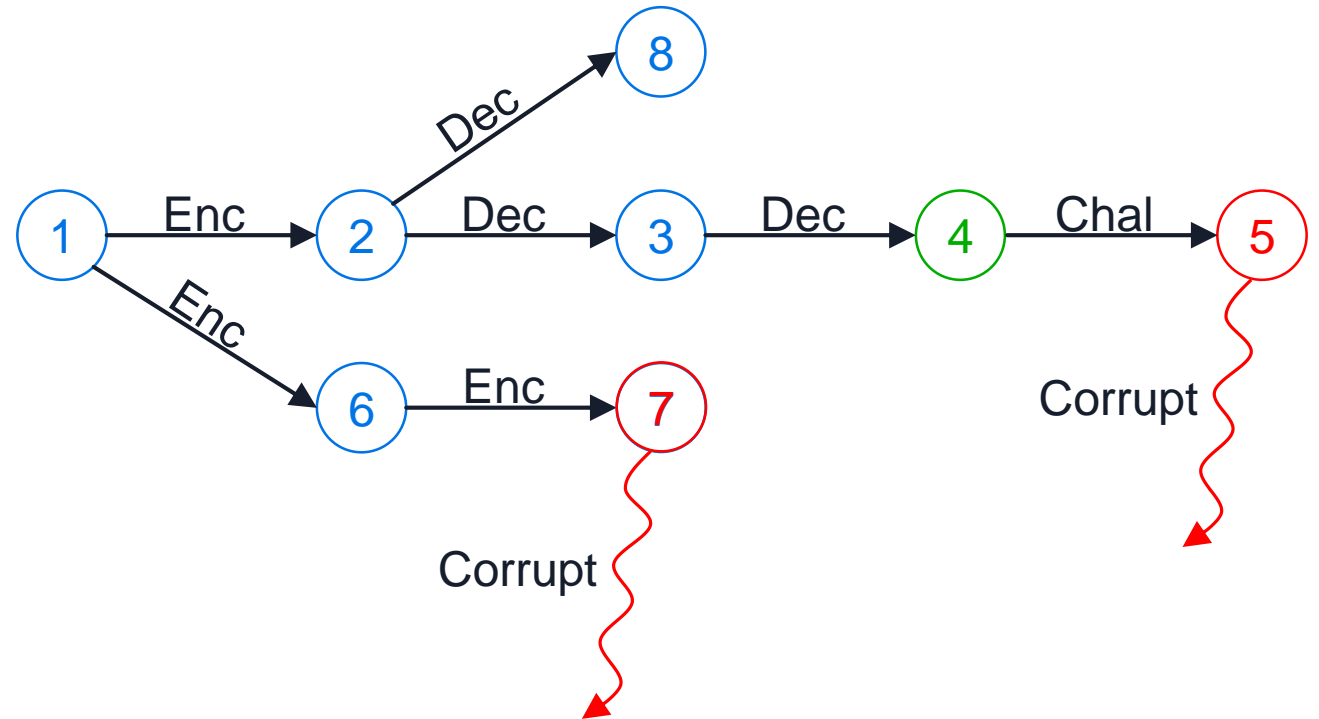
$(K, sk') \leftarrow Dec(sk, c)$

$0/1 \leftarrow Verify(pk, pk', mt')$

IND-CCA-Members [this work]

what do we need for Secure Group Messaging?

- “Enc-forks”
- “Dec-forks”
- multiple corruptions
- challenge most nodes



UKEM Syntax and Security



Extended Syntax for Joiner Security [\[this work\]](#)

$(pk1, sk1) \leftarrow KGen()$

$(K, c, pk', mt') \leftarrow Enc(pk)$

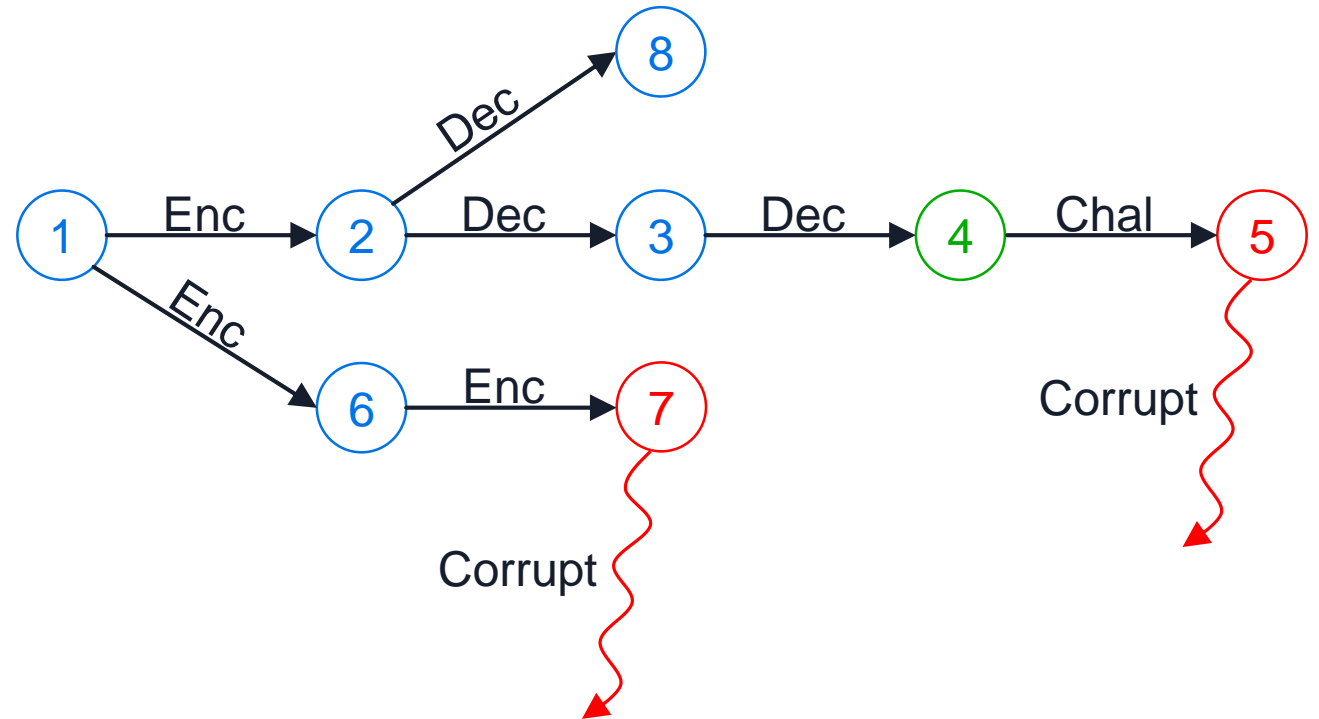
$(K, sk') \leftarrow Dec(sk, c)$

$0/1 \leftarrow Verify(pk, pk', mt')$

Motivation

Alice joins mid-session

- everyone else is offline
- keys are at epoch 8
- Alice trusts that epoch 1 is honest
- how can Alice trust pk at epoch 8?



UKEM Syntax and Security



Extended Syntax for Joiner Security [\[this work\]](#)

$(pk1, sk1, jt1) \leftarrow KGen()$

$(K, c, pk', mt', jt') \leftarrow Enc(pk, jt)$

$(K, sk') \leftarrow Dec(sk, c)$

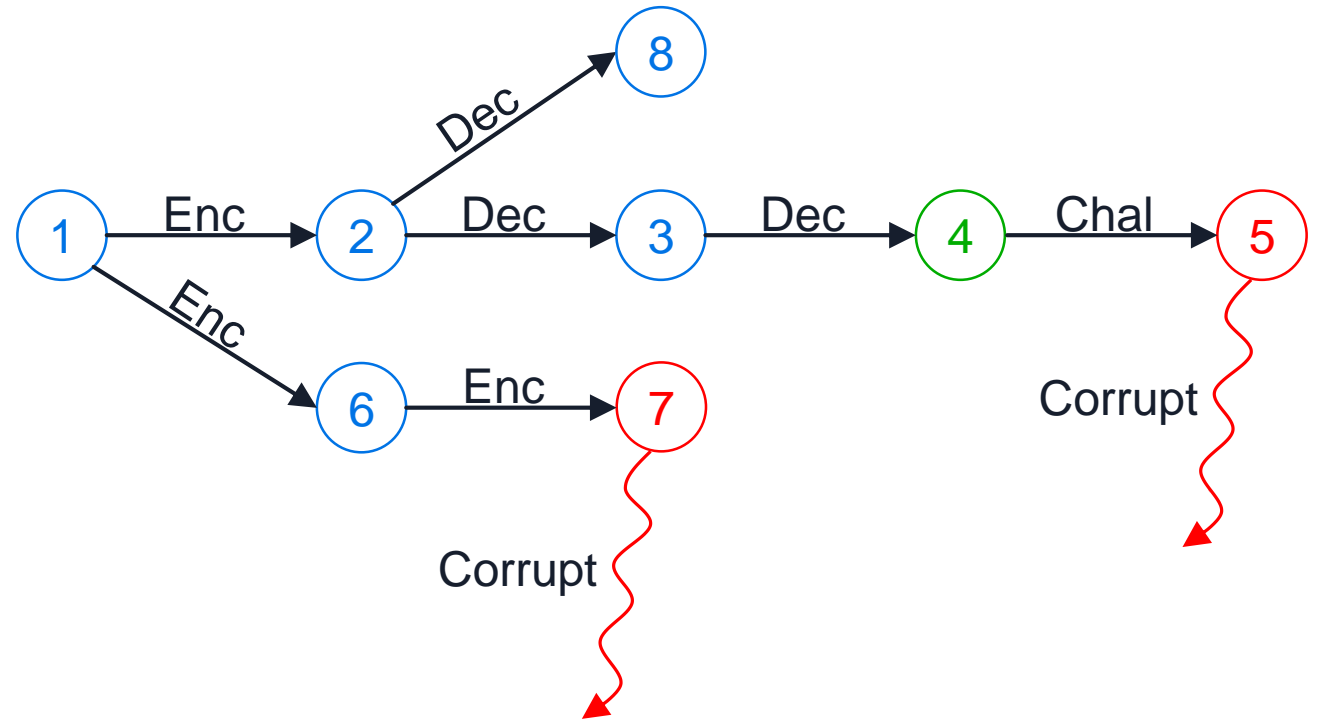
$0/1 \leftarrow Verify(pk, pk', mt')$

$0/1 \leftarrow VerifyJT(pk1, pk', jt')$

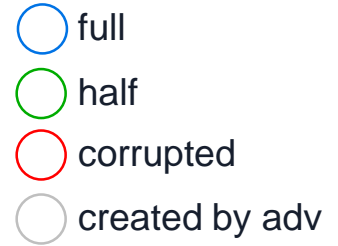
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UKEM Syntax and Security



Extended Syntax for Joiner Security [this work]

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$(K, c, pk', mt', jt') \leftarrow Enc(pk, jt)$

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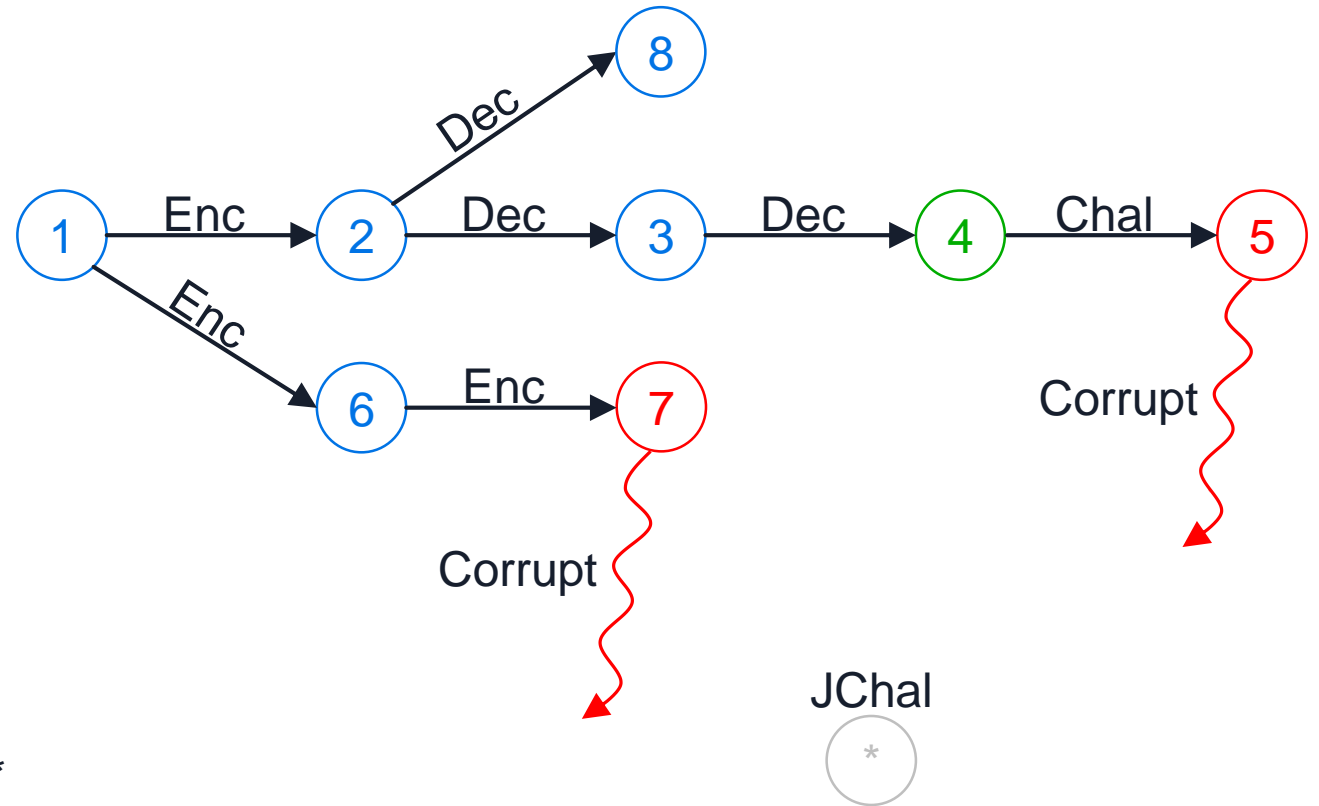
$0/1 \leftarrow Verify(pk, pk', mt')$

$0/1 \leftarrow VerifyJT(pk1, pk', jt')$

IND-CCA-Joiners [this work]

add oracle $JChal(pk', jt')$

- pk', mt' arbitrary
- if $VerifyJT$ succeeds
 - runs $Enc(pk', jt')$ and returns c^* , real or random K^* (and other Enc outputs)



Construction

DHIES-KEM

assumes : group G with generator g

key pair : $(x, X = g^x)$

Enc(X)

$r \leftarrow \$$

$R := g^r$

$K := H_1(X^r, X, R)$

return K, R

Construction

DHIES-KEM with **key update**

assumes : group G with generator g

key pair : $(x, X = g^x)$

Enc(X)

$r \leftarrow \$$

$R := g^r$

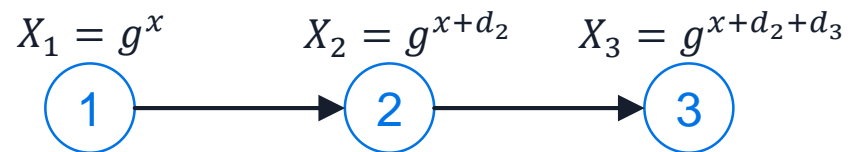
$K := H_1(X^r, X, R)$

$d := H_2(X^r, X, R)$

$X' := X \cdot g^d (= g^{x+d})$

return K, R, X'

using H_2 allows to deal with adaptive corruptions in the proof
→ with linear security loss!



Construction

DHIES-KEM with key update and **member tags**

assumes : group G with generator g

key pair : $(x, X = g^x)$

Enc(X)

$r \leftarrow \$$

$R := g^r$

$K := H_1(X^r, X, R)$

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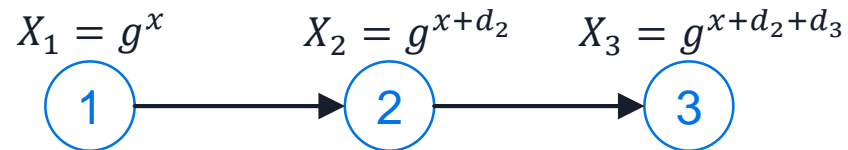
$X' := X \cdot g^d \quad (= g^{x+d})$

$\tau \leftarrow \text{Schnorr-ZKPoK}(d)$

return K, R, X', τ

strong simulation (multi-) extractable

τ_3 proves knowledge of d_3



Construction

DHIES-KEM with key update, member tags and **joiner tags**

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key pair : $(x, X = g^x)$

Enc(X)

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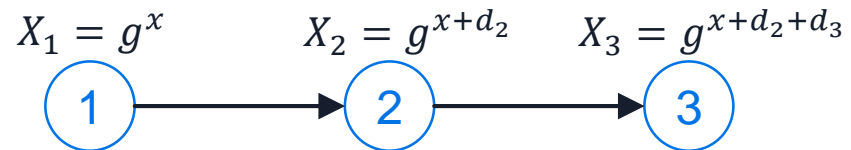
$d := H_2(X^r, X, R)$

$X' := X \cdot g^d \quad (= g^{x+d})$

$\tau' \leftarrow \text{Schnorr-ZKPoK}(d)$

$\pi' \leftarrow \text{ZKPoK}(d_2 + d_3 + \dots + d)$

return K, R, X', τ, π'



Construction

DHIES-KEM with key update, member tags and **joiner tags**

assumes : group G with generator g

key pair : $(x, X = g^x)$

Enc (X, π)

$r \leftarrow \$$

$R := g^r$

$K := H_1(X^r, X, R)$

$d := H_2(X^r, X, R)$

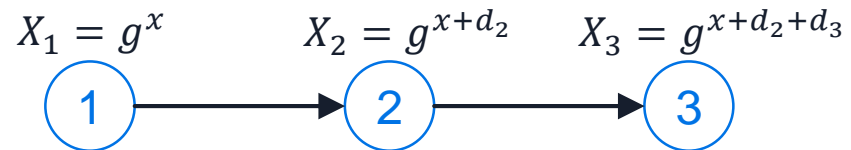
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return K, R, X', τ, π'

constant-size malleable ZKPoK



Summary

- 2 new notions and constructions of Updatable KEMs

| UKEM security | Assumption | pk + ctx | joiner tag |
|-----------------|--------------------------|------------|------------|
| IND-CCA-Members | DH (ROM) | 0.24 kB | – |
| IND-CCA-Joiners | DH + pairing (ROM + AGM) | 0.24 kB | 0.01 kB |

- New tool : malleable ZKPoK